

IBM Research – Tokyo

Workload Characterization and Optimization of TPC-H Queries on Apache Spark

Tatsuhiro Chiba and Tamiya Onodera

IBM Research - Tokyo

April. 17-19, 2016 | IEEE ISPASS 2016 @ Uppsala, Sweden

© 2016 IBM Corporation

Overview

- Introduction
 - Motivation
 - Goal and Result
- Workload Characterization
 - How Spark and Spark SQL work
 - Environments
 - Application level analysis
 - System level analysis
 - GC analysis
 - PMU analysis
- Problem Assessments and Approach
- Result
- Summary and Future Work

Motivation

Apache Spark is an in-memory data processing framework, runs on JVM Spark executes Hadoop similar workloads, but *optimization points are not same*

- Complexity to find a fundamental Bottlenecks
 - I/O bottleneck \rightarrow CPU, Memory, Network bottleneck
 - JVM handles many worker threads

→ What is a best practice to achieve high performance?

- Managing large Java heap causes high GC overhead
 - Keeps as much data in memory as possible
 - Generates short-lived immutable Java objects

→ How we can reduce garbage collection overhead?

- From Scale-out to Scale-up
 - Utilizes many worker threads w/ SMT on multiple NUMA nodes
 - Need to know micro architectural efforts

→ How we can exploit underlying Hardware features?



3



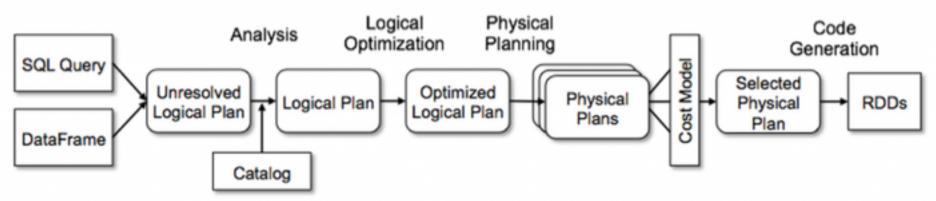
Goal and Result

- Goal
 - To characterize Spark Performance through SQL workloads
 - To find optimization best practice for Spark
 - To investigate potential problem in current Spark and JVM for scaling up
- Result
 - Achieved up to 30% improvement by reducing GC overhead
 - Achieved up to 10% improvement by utilizing more SMT
 - Achieved up to 3% improvement by NUMA awareness
 - Achieved 30 40% improvement on average with applying all optimization



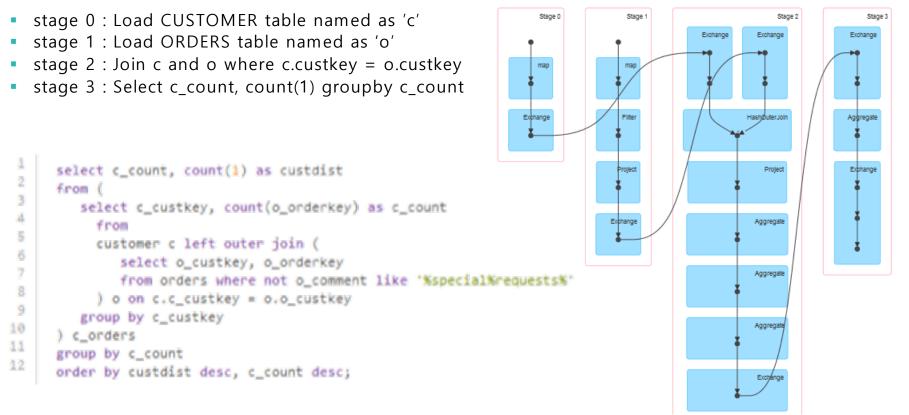
How Spark and Spark SQL work

- Spark
 - Job is described as a data transformation chain and divided into multiple stages
 - Each stage includes multiple tasks
 - Tasks are concurrently proceeded by worker threads on JVMs
 - Data shuffling occurs between stages
- Spark SQL
 - Catalyst, a query optimization framework for Spark, generates an optimized code
 - It has a compatibility for HIVE query



Cited from Michael et al., Spark SQL: Relational Data Processing in Spark , SIGMOD'15

How SQL code translates into Spark - TPC-H Q13



Completed Stages (4)

Stage Id	Description	Submitted	Duration	Tasks: Succeeded/Total	Input	Output	Shuffle Read	Shuffle Write
3	select c_count, count(1) as custdist from (select c_custkey, count(o_orderkey) as c_count from Spark JDBC Server Query +details	2015/06/24 04:35:34	0.4 s	200/200			693.4 KB	
2	select c_count, count(1) as custdist from (select c_custkey, count(o_orderkey) as c_count from Spark JDBC Server Query +details	2015/06/24 04:34:47	47 s	200/200			15.0 GB	693.5 KB
1	select c_count, count(1) as custdist from (select c_custkey, count(o_orderkey) as c_count from Spark JDBC Server Query +details	2015/06/24 04:31:48	3.0 min	1348/1348	33.3 GB			14.2 GB
0	select c_count, count(1) as custdist from (select c_custkey, count(o_orderkey) as c_count from Spark JDBC Server Query +details	2015/06/24 04:31:48	12 s	185/185	400.1 MB			817.6 MB

Machine & Software Spec and Spark Settings

Processor # Core		SMT	SMT Memory		
POWER824 cores3.30 GHz * 2(2 sockets * 12 cores)		8 (total 192 hardware threads)	1TB	Ubuntu 14.10 (kernel 3.16.0-31)	
	software	version			
	Spark	1.4.1			
Ha	adoop (HDFS)	2.6.0			
	Java	1.8.0 (IBM J9 VM SR1 FP10)			
	Scala	2.10.4			

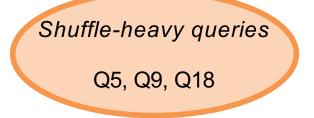
- Baseline Spark settings
 - # of Executor JVMs: 1
 - # of worker threads: 48
 - Executor heap size: 192GB (nursery = 48g, tenure = 144g)
- Other picked up Spark configurations (spark-defaults.conf)
 - spark.suffle.compress = true
 - spark.sql.parquet.compression.codec = Snappy
 - spark.sql.parquet.fileterPushdown = true

Workload Characterization – Spark job level

* Picked up several queries

Query	SQL Characteristics	Converted Spark Operation (# of stages)	Input (total, GB)	Shuffle (total, GB)	Stages / Tasks	Time (sec)	
Q1	1 GroupBy Load 1 Table	1 Load 1 Aggregate	4.8	0.002	2 / 793	48.7	
Q3	1 GroupBy, 2 Join Load 3 Table	3 Load 2 HashJoin, 1 Aogregate	7.3	5.0	6 / 1345	64.6	
Q5	1 GroupBy, 5 Join Load 6 Table	3 Load, 3 HashJoin 1 BcastJoin, 1 Aggregate	8.8	14.1	8 / 1547	125	
Q6	1 Select, 1 Where Load 1 Table	1 Load 1 Aggregate	4.8	0	2/594	15.1	
Q9	1 GroupBy, 5 Join Load 6 Table	4 Load, 4 HashJoin 1 BcastJoin, 1 Aggregate	11.8	34.4	10 / 1838	370	
Q18	3 Join, 1 UnionAll Load 3 Table	6 Load, 3 HashJoin 1 Union, 1 Limit	7.7	13.8	11 / 3725	202	
Q19	3 Join, 1UnionAll Load 2 Table	6 Load, 3 HashJoin 1 Union, 1 Aggregate	19.8	0.4	8 / 2437	80.8	

Shuffle-light queries	
Q1, Q6, Q19	



Workload Characterization and Optimization for TPC-H Queries on Apache Spark

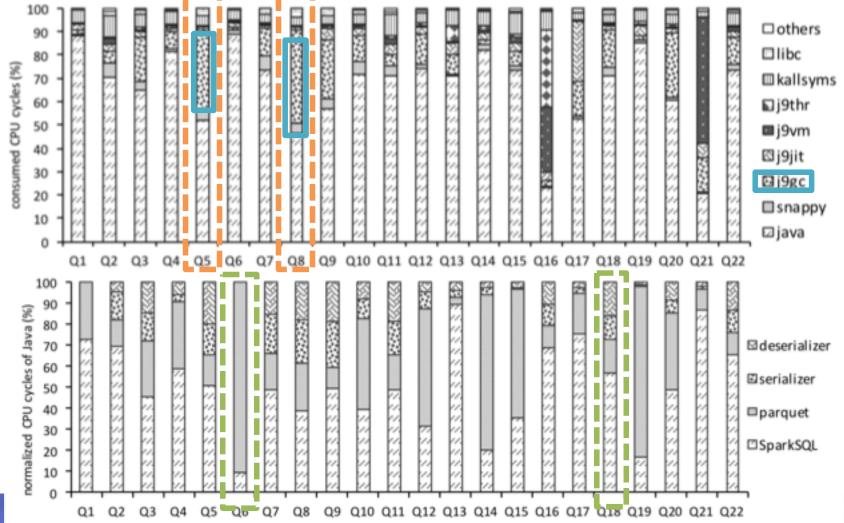


IBM Research – Tokyo

9

Workload Characterization – oprofile

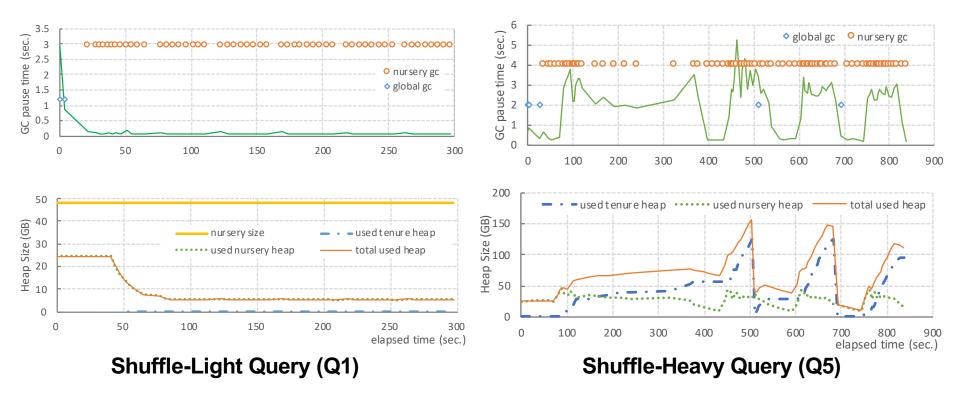
- Shuffle-heavy queries (e.g. Q5 and Q8) : over 30% cycles are spent in GC Shuffle-light queries (e.g. Q1 and Q6) : low SerDes cost
- Unexpected JIT and Thread Spin Lock overheads exists in Q16 and Q21



Workload Characterization – garbage collection

- GC
 - Many nursery GC ab
 - Small Pause time (0.02 sec)
 - Few global GC
- Java Heap
 - Low usage level (within nursery space)

- GC
 - Many nursery GC
 - Big pause time (3 5 sec)
 - Global GC while execution
- Java Heap
 - Objects are flowed into tenure space



Workload Characterization – PMU profile

- Approach
 - Observed performance counters by perf
 - Categorized them based on the CPI breakdown model [*]
- Result
 - Backend stalls are quite big
 - Lots of L3 miss which comes from distant memory access
 - CPU Migration occurs frequently

counters	Q1	Q5
CPU cycles	6.8×10^{12}	2.2×10^{13}
stalled-cycles-frontend	2.1×10^{11} (3.20%)	6.2×10^{11} (2.76%)
stalled-cycles-backend	3.3×10^{12} (49.0%)	1.3×10^{13} (59.1%)
instructions	$7.0 imes 10^{12}$	$1.5 imes 10^{13}$
IPC	1.03	0.67
context-switches	407K	440K
cpu-migrations	11K	26K
page-faults	308K	1045K

[*] https://www.ibm.com/support/knowledgecenter/linuxonibm/liaal/iplsdkcpieventspower8.htm

Workload Characterization and Optimization for TPC-H Queries on Apache Spark





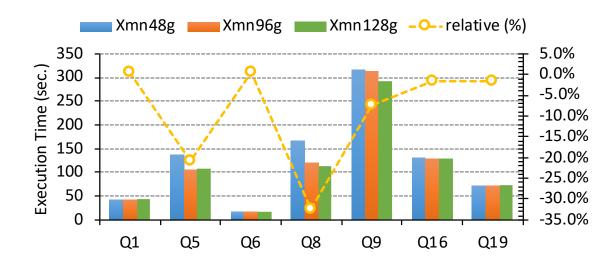
Problem Assessments and Optimization Strategies

- How we can reduce GC overhead?
 - 1). Heap sizing
 - -2). JVM Option tuning
 - -3). Changing # of Spark Executor JVMs
 - -4). GC algorithm tuning
- How we can reduce backend stall cycles?
 - 1). NUMA awareness
 - -2). Adding more hardware threads to Executor JVMs
 - 3). Changing SMT level (SMT2, SMT4, SMT8)

Efforts of Heap Space Sizing

Heap sizing efforts

- Bigger nursery space achieves up to 30% improvement
- Small tenure space may be harmful
 - Run out of tenure space by caching RDDs in memory
 - Leaked objects from nursery space



TPC-H Q9 Case

Nursery Space (-Xmn)	Execution Time (sec)	GC ratio (%)	Nursery GC Avg. pause time	Nursery GC	Global GC
48g (default)	316 s	20 %	2.1 s	39	1
96g	310 s	18 %	3.4 s	22	1
144g	292 s	14 %	3.6 s	14	0

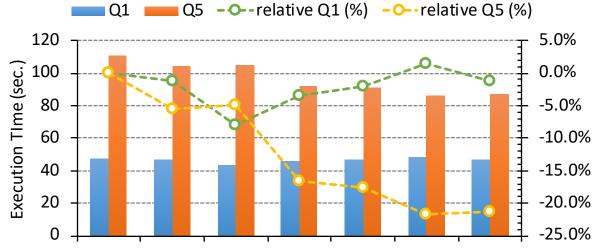
Workload Characterization and Optimization for TPC-H Queries on Apache Spark

© 2016 IBM Corporation

IBM

Efforts of Other JVM Options

- JVM options
 - Monitor threads tuning
 - # of GC threads tuning
 - Java thread tuning
 - JIT tuning, etc.
- Result
 - Improved over 20%



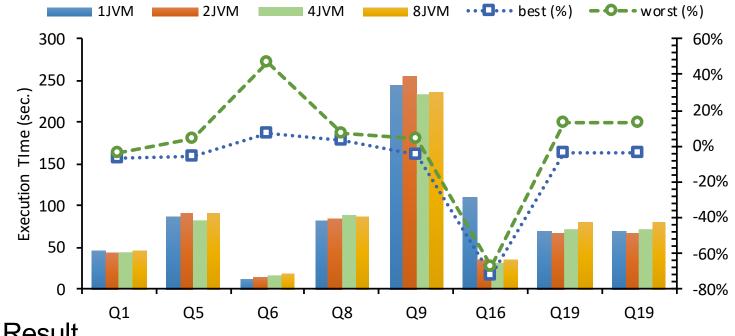
option 1option 2option 3option 4option 5option 6option 7

	#	spark.executor.extraJavaOptions	
GC Threads \rightarrow		-Xgcthreads48 -Xmn96g -Xdump:system:none -Xdump:heap:none	
		-Xtrace:none	
Monitor Threads 🔿	2	-Xgcthreads48 -Xmn96g -Xdump:system:none -Xdump:heap:none	
		-Xnoloa -Xtrace:none	
Disable Large Object Area →	3	-Xgcthreads48 -Xmn96g -Xdump:system:none -Xdump:heap:none	
C <i>j</i>		-XlockReservation -Xnoloa -Xtrace:none	
Reduce thread lock cost →	4	-Xgcthreads48 -Xmn96g -Xdump:system:none -Xdump:heap:none	
		 -Xnocompactgc -XlockReservation -Xnoloa -Xtrace:none 	
Stop compaction \rightarrow	5	-Xgcthreads48 -Xmn96g -Xdump:system:none -Xdump:heap:none	
		-XX:-RuntimeInstrumentation	
Stop JIT feature 🔿		-Xnocompactgc -XlockReservation -Xnoloa -Xtrace:none	
	6	-Xgcthreads48 -Xmn96g -Xdump:system:none -Xdump:heap:none	
		-Xdisableexplicitge -XX:-RuntimeInstrumentation	
Stop System.gc() →		-Xnocompactgc -XlockReservation -Xnoloa -Xtrace:none	
	7	-Xgcthreads48 -Xmn96g -Xdump:system:none -Xdump:heap:none	

Workload Characterization and Optimization for TPC-H Queries on Apache Spark



Efforts of changing JVM Counts

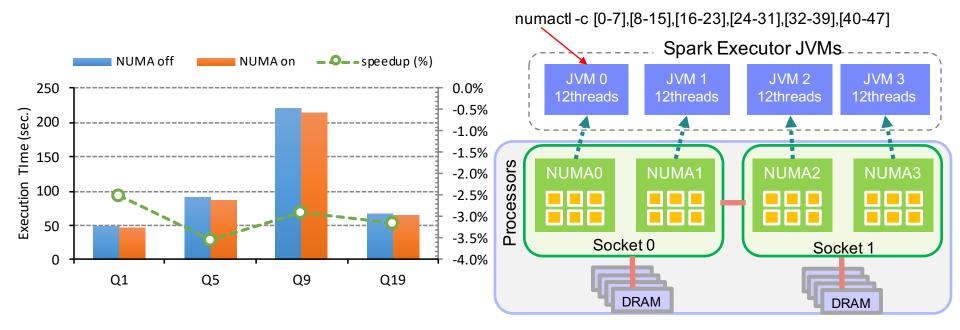


Result

- Up to 70% improvement in Q16 (by avoiding unexpected threads lock activity)
- Reduced heavy GC overhead
- Has a drawback a little for shuffle-light queries
- Using 1 JVM frequently occurs task execution failure than 4 JVMs



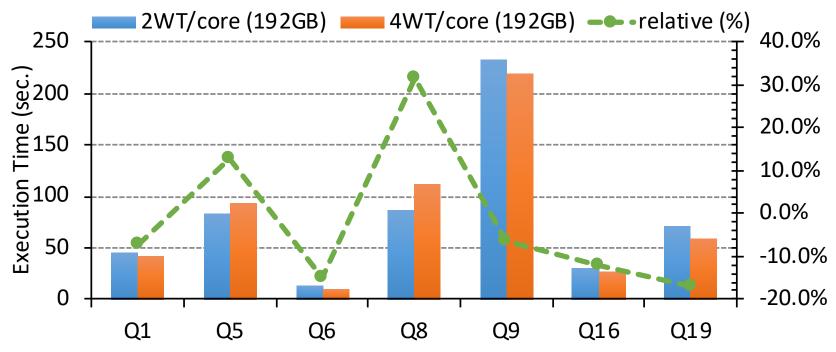
Efforts of NUMA aware process affinity



- Setting NUMA aware process affinity to each Executor JVM helps to speedup
 - By reducing scheduling overhead
 - By reducing cache miss and stall cycles
- Result
 - Achieved 3 3.5 % improvement in all benchmarks without any bad effects



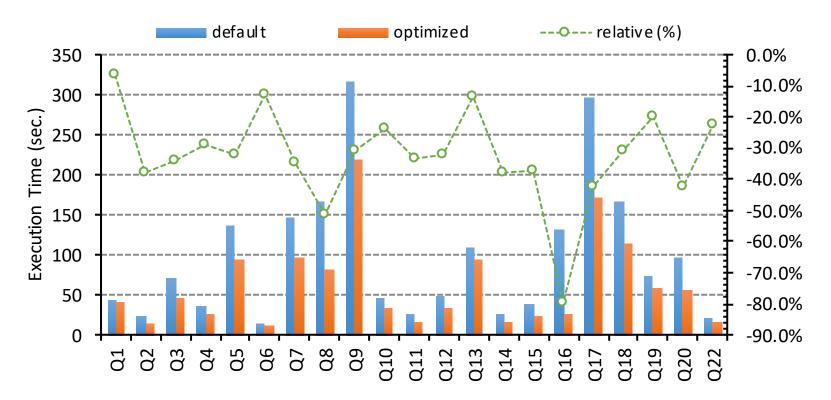
Efforts of Increasing worker threads



- Settings
 - 2WT/core : handles 12 worker threads on 6 cores (in total, 48 worker threads)
 - 4WT/core : handles 24 worker threads on 6 cores (in total, 96 worker threads)
- Result
 - Some queries gain over 10% improvement regardless of shuffle data size
 - Q5 and Q8 had a drawback



Summary of applying all optimizations



- Shuffle-light: 10 20% improvement
- Shuffle-heavy: 30 40% improvement
- Eliminated unexpected JVM behavior in Q16 and Q21
 - Q21 took 543 sec, which took over 3000 sec before tuning



Summary and Future Works

- Summary
 - Reduced GC overhead from 30% to 10% or less by heap sizing, JVM counts, and JVM options
 - Reduced distant memory access from 66.5% to 58.9% by NUMA awareness
 - In summary, achieved 30 40 % improvement on average
 - All experiment codes are available at <u>https://github.com/tatsuhirochiba/tpch-on-spark</u>
- Future works
 - Comparison between x86_64 and POWER8
 - Other Spark workloads